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A NOTE ON OMITTING THE REPLACEMENT SCHEMA

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In [1] Heath considers a formalisation of primitive recursive arithmetic similar to that given in Goodstein [2], in which the replacement schema (Goodstein's $Sb_2$) is deduced from special cases of itself, using a double recursive uniqueness rule. The deduction of $Sb_2$ given in [1] is, however, incomplete. This is rectified in the present note. The special cases of $Sb_2$ taken by Heath are:

(i) $A = B \vdash SA = SB$
(ii) $A = B \vdash x + A = x + B$
(iii) $A = B \vdash A + x = B + x$
(iv) $A = B \vdash \frac{x}{A} = \frac{B}{x}$
(v) $A = B \vdash A \times A = B \times x$

Remark: In fact either (ii) or (iii) can be omitted since $x + y = y + x$ can be proved without using (ii) or (iii) and then one can be derived from the other.

In order to derive the full $Sb_2$, i.e., $A = B \vdash f(A) = f(B)$, for any primitive recursive function $f$, it is necessary to show that the substitution theorem, $x = y \rightarrow f(x) = f(y)$, persists under definition by a primitive recursive schema. Heath shows that it persists under the recursion without parameter, which I shall call $R$,

$f(0) = (0)$,
$f(Sx) = g(x, f(x))$,

i.e., that from $x = y$ & $w = z \rightarrow g(x, w) = g(y, z)$ we can deduce $x = y \rightarrow f(x) = f(y)$. He then quotes a theorem of R. M. Robinson that all primitive recursive functions are generated from 0, $x$, $Sx$, $x + y$ and $x \div y$ by substitution and the recursion $R$. To complete the proof it would be sufficient to show that Robinson's reduction of primitive recursion can be carried out in the restricted primitive recursive arithmetic (i.e., without full $Sb_2$). This would involve defining the pairing functions $J(x, y)$, $K(x)$ and $L(x)$ given by Robinson, deriving their main properties, e.g. $L(Sx) \neq 0 \rightarrow K(Sx) = K(x)$ & $L(Sx) = S(L(x)$, and checking that the substitution theorem is satisfied by them. This part was omitted by Heath, and it is not clear that this programme could be carried out.

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However it is fairly easy to check that the substitution theorem persists under full recursion, by a simple adaptation of Heath's proof for the recursion scheme \( R \), as the following theorem shows.

**Theorem** Suppose \( f \) is defined by primitive recursion from \( h \) and \( g \), i.e.,

\[
\begin{align*}
f(u_0, \ldots, u_n, 0) &= h(u_0, \ldots, u_n) \\
f(u_0, \ldots, u_n, Sx) &= g(u_0, \ldots, u_n, x, f(u_0, \ldots, u_n, x))
\end{align*}
\]

and the substitution theorem has already been proved for \( h \) and \( g \), i.e.,

\[
\begin{align*}
u_0 = v_0 &\land \ldots \land u_n = v_n \rightarrow h(u_0, \ldots, u_n) = h(v_0, \ldots, v_n) \\
u_0 = v_0 &\land \ldots \land u_n = v_n \rightarrow g(u_0, \ldots, u_n, x, f(u_0, \ldots, u_n, x)) = g(v_0, \ldots, v_n, x, f(v_0, \ldots, v_n, x))
\end{align*}
\]

Then the substitution theorem holds for \( f \), i.e.,

\[
u_0 = v_0 &\land \ldots \land u_n = v_n \rightarrow f(u_0, \ldots, u_n, 0) = f(v_0, \ldots, v_n, 0)
\]

**Proof**

**Lemma I** \( u_0 = v_0 & \ldots & u_n = v_n \rightarrow f(u_0, \ldots, u_n, x) = f(v_0, \ldots, v_n, x) \)

By induction on \( x \), prove the basis

\[
u_0 = v_0 &\land \ldots \land u_n = v_n \rightarrow f(u_0, \ldots, u_n, 0) = f(v_0, \ldots, v_n, 0)
\]

by hypotheses (a) and (c)

and the step

\[
u_0 = v_0 &\land \ldots \land u_n = v_n & (u_0 = v_0 & \ldots & u_n = v_n \rightarrow f(u_0, \ldots, u_n, x) = f(v_0, \ldots, v_n, x)) \rightarrow f(u_0, \ldots, u_n, Sx) = f(v_0, \ldots, v_n, Sx)
\]

by hypotheses (b) and (d).

**Lemma II** \( x = y \rightarrow f(u_0, \ldots, u_n, x) = f(u_0, \ldots, u_n, y) \)

By double induction on \( x \) and \( y \), prove

\[
x = 0 \rightarrow f(u_0, \ldots, u_n, x) = f(u_0, \ldots, u_n, 0)
\]

and

\[
0 = y \rightarrow f(u_0, \ldots, u_n, 0) = f(u_0, \ldots, u_n, y)
\]

by schema \( F \) on \( x \) and \( y \) respectively. Then use the deduction theorem to prove

\[
(x = y \rightarrow f(u_0, \ldots, u_n, x) = f(u_0, \ldots, u_n, y)) \rightarrow (Sx = Sy \rightarrow f(u_0, \ldots, u_n, Sx) = f(u_0, \ldots, u_n, Sy))
\]

Assume \( x = y \rightarrow f(u_0, \ldots, u_n, x) = f(u_0, \ldots, u_n, y) \) and \( Sx = Sy \) and without using \( Sb \) on any of the variables \( u_0, \ldots, u_n, x, y \), deduce, in turn,

\[
x = y \\rightarrow f(u_0, \ldots, u_n, x) = f(u_0, \ldots, u_n, y)
\]

by modus ponens

\[
g(u_0, \ldots, u_n, x, f(u_0, \ldots, u_n, x)) = g(u_0, \ldots, u_n, y, f(u_0, \ldots, u_n, y))
\]

by hypothesis (d).
Therefore

\[ f(u_0, \ldots, u_n, Sx) = f(u_0, \ldots, u_n,Sy) \]

by hypothesis (b).

The theorem follows from Lemmas I and II.

**REFERENCES**


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