Incremental Evaluation of Schema-directed XML Publishing

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Abstract

When large XML documents published from a database are maintained externally, it is inefficient to repeatedly recompute them when the database is updated. Vastly preferable is incremental update, as common for views stored in a data warehouse. However, to support schema-directed publishing, there may be no simple query that defines the mapping from the database to the external document. To meet the need for efficient incremental update, this paper studies two approaches for incremental evaluation of ATGs [4], a formalism for schema-directed XML publishing. The reduction approach seeks to push as much work as possible to the underlying DBMS. It is based on a relational encoding of XML trees and a nontrivial translation of ATGs to SQL 99 queries with recursion. However, a weakness of this approach is that it relies on high-end DBMS features rather than the lowest common denominator. In contrast, the bud-cut approach pushes only simple queries to the DBMS and performs the bulk of the work in middleware. It capitalizes on the tree-structure of XML views to minimize unnecessary recomputations and leverage optimization techniques developed for XML publishing. While implementation of the reduction approach is not yet in the reach of commercial systems, we have implemented the bud-cut approach and experimentally evaluated its performance compared to recomputation.

1. Introduction

XML publishing by middleware [11, 8, 16] or with direct DBMS support [7] has been well studied, and techniques from this work are rapidly being introduced into commercial products [25, 28]. In some applications, small portions of a database are extracted into “disposable” XML documents, for example the messages needed to execute or respond to requests using a web-services protocol. However, in many applications including mediation, archiving and web site management, large XML documents may need to be exported. In this case, the cached XML document can obviously be seen as a view of the database instance. For all the reasons that efficient incremental maintenance of views in the database (see, e.g., [14]) is important, it may also make sense to incrementally update published XML documents, even when they are externally cached by a middleware system. However, to

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\[ I : R \rightarrow \alpha \]
\[ T : D \]
\[ I \oplus \Delta \]
\[ \alpha \]
\[ \Delta T \]
\[ T \ominus \Delta T : D \]

Figure 1: Incremental evaluation

our knowledge, no previous work has considered incremental update of XML documents published from relational data.

In this paper, we consider the particular case of incremental update of XML documents created by schema-directed XML publishing middleware. The idea of schema-directed publishing is to extract data from a relational database and construct an XML view that \textit{conforms to a predefined schema}, in this case a DTD. The need for this is evident in practice: enterprises typically agree on a common schema for data exchange; thus, a specification for publishing views that ensures schema conformance is an obvious benefit for developers. In response to this need, the ATG (Attribute Translation Grammar) formalism has been developed [4]. An ATG is a mapping \( \sigma : R \rightarrow D \) associated with a relational schema \( R \) and a predefined (possibly recursive) DTD \( D \). Given an instance \( I \) of \( R \), \( \sigma \) computes an XML view \( T = \sigma(I) \) such that \( T \) conforms to \( D \). To accomplish this transformation, \( D \) is treated as a set of productions. Each production is annotated by \( \sigma \) with a set of \textit{semantic rules}, and a single \textit{semantic attribute} is defined for each element type. These rules govern the production of child elements for each element based on the data present in the attribute.

To discuss incremental update, we assume that a set of changes to an XML tree \( T \) can be encapsulated as \( \Delta T \), and that an operator \( \oplus \) represents the application of these updates. Given these assumptions, the \textit{incremental evaluation problem} for ATGs can be stated as follows: given an ATG \( \sigma : R \rightarrow D \), a relational instance \( I \) of \( R \), the XML view \( T = \sigma(I) \), and changes \( \Delta T \) to \( I \), compute XML changes \( \Delta T \) to \( T \) such that \( T \oplus \Delta T = \sigma(I \oplus \Delta I) \). These relationships are illustrated in Fig. 1. Note further that, if \( \sigma \) is correct, \( T \ominus \Delta T \) is trivially guaranteed to conform to the predefined DTD \( D \). In contrast to recomputing the new view from scratch, incremental evaluation of ATGs can, in principle, improve performance substantially by applying only the changes \( \Delta T \) to the old view \( T \). To realize the improvement requires a) an efficient algorithm that computes the XML changes \( \Delta T \) in response to relational changes \( \Delta I \), and b) an efficient implementation of the tree-update operator \( \oplus \).

Our first contribution is a \textit{reduction framework} for incremental maintenance of XML views defined by ATGs. For a given ATG definition \( \sigma \), the framework consists of a) a set of virtual relations associated with the semantic attributes of \( \sigma \), b) queries defining the recursive relationship of these virtual relations to encode ATG-produced XML documents c) a mapping based on these recursive queries from ATGs to SQL 99 queries with linear recursion [24], and d) mechanisms for computing XML changes \( \Delta T \) from relational changes \( \Delta I \)
and for updating the external view $T$ with $\Delta T$.

Previous work [11, 8] for XML publishing has developed mappings from XML view queries to SQL queries that compute root-to-leaf paths in a relational encoding of the XML tree. These techniques however, cannot be directly applied to generating or updating AVG-generated documents due to a lack of support for recursion. In the context of XML shredding, the authors of [17] showed that linear recursion of SQL 99 is sufficient to support XPath queries over shredded XML data, even when the shredding schema is recursive, and suggested that this approach could be applied to publishing queries. The reduction framework extends the prior work by showing the connection between the recursive XML views needed for schema-directed publishing and SQL 99 views.

Much as [11, 8] seek to push XML publishing work to the DBMS, a primary goal of the reduction framework is to push incremental work to the DBMS, thereby taking advantage of sophisticated capabilities for query optimization, execution, and incremental view update. However, three practical issues complicate the use of DBMS resources in support of incremental work. The first issue involves the DBMS features required for the reduction approach. Middleware that seeks to work with a wide variety of products supporting the ODBC interface must take a "lowest common denominator" approach to the functionality required from the DBMS. But SQL 99 recursion and incremental update of views are separate, advanced features of only the most sophisticated commercial products; furthermore the reduction approach also depends on incremental update of materialized views defined using SQL 99 recursion; in other words, a combination of both features. Second, to effectively push down the work required to incrementally update an external view, one must have access not only to a materialized view, but more importantly, to a stream of updates to that view. One way to obtain this functionality would be to define triggers on the materialized view, but this is not supported by at least one commercial DBMS with materialized view support, and needs not be supported in general for the DBMS to function well. Finally, if the publishing queries are even mildly complex, the combined recursive queries may become extremely complex. As a result, they may not be effectively optimized by all platforms supporting with...for the same reasons that not all DBMS platforms can effectively optimize complex non-recursive publishing queries [11].

In response to this, our second contribution is to propose an alternative approach, referred to as the "bud-cut approach" to incremental ATG evaluation, that requires less sophistication from the DBMS. Further, we develop certain optimization techniques which capitalize on the tree-structure of XML views. The bud-cut mechanism propagates relational changes to XML in three phases: generation, completion and garbage collection. The bud-cut generation phase determines the impact of $\Delta I$ on existing parent-child relations in the XML view, i.e., insertions (buds) and deletions (cubs), by evaluating a fixed number of parameterized SQL queries. Following this, the bud completion phase iteratively computes newly inserted subtrees top-down by pushing SQL queries to the relational engine. Finally, deleted subtrees are removed by a garbage collection process.

The bud-cut approach has several properties. a) $T$ can be updated in parallel with ongoing computation of $\Delta T$ during bud completion. b) It minimizes unnecessary recomputations via a caching strategy not considered in prior work for maintaining recursive views such that each new subtree in the XML view is computed at most once no matter how many times it occurs in the XML view, and furthermore, the computation maximally reuses subtrees of the old XML view. c) It incorporates optimization techniques that have proved effective in XML publishing but are not supported by DBMSs, e.g., query merging [11]. d) Since the tree is computed level-by-level in this phase, it is possible to return partial results to a user navigating the tree while computation is ongoing, and such computation can even be deferred according to a lazy evaluation strategy [7]. e) It does not require materialization of the view in the DBMS. Finally, of course, the bud-cut approach does not require the DBMS to support either SQL 99 or incremental view updates.

We have implemented the bud-cut approach. We use the implementation to investigate the impact of $I$ and $\Delta I$ on the performance of incremental update, as well as the improvement obtained over full recomputation for small updates. Further, we investigate the impact on performance of the subtree reuse optimization mentioned above, and find that its impact is greatest when only a portion of $I$ is published in $T$, and when there is a moderate degree of locality in the updates appearing in $\Delta I$. Unfortunately but not surprisingly, the reduction approach is unrealistic since current commercial relational systems do not implement incremental update of recursive queries.

The algorithms and the bud-cut mechanism can be extended to accommodate multiple data sources, i.e., for XML integration studied by [3], which is a generalization of ATG's by supporting multi-source SQL queries and XML constraints. They can also be used for incremental maintenance of XML views generated by other systems, such as [11, 8].

**Organisation.** Section 2 reviews ATG. Section 3 describes data structures for external XML views. Section 4 provides the reduction approach, followed by the bud-cut approach in Sect. 5. Section 6 presents experimental results. Section 7 addresses related work and Sect. 8 concludes the paper.

**2. Background**

In this section we first introduce a running XML publishing example used in the rest of the paper. We then review DTDs and present a refinement of ATGs as defined by [4].

**Example 2.1:** Consider a registrar database specified by the relational schema $R_0$ below (with keys underlined):

- course(cno, title, dept),
- project(cno, title, dept),
- student(name),
- enroll(cno, cno),
- prereq(cno, cno).

The database maintains student data, enrollment records, course data classified into regular courses and projects, and a relation prereq which gives the prerequisite hierarchy of courses where a tuple $(c_1, c_2)$ in prereq indicates that $c_2$ is a prerequisite of $c_1$.

The office of registrar maintains an XML view for the CS department, which contains data of CS courses registration, extracted from the registrar database. The view is required to conform to the DTD $D_0$ below (the definition of elements whose type is PCDATA is omitted):

```xml
<ELEMENT db (course)>
<ELEMENT course (cno, title, type, prereq, takenby)>
<ELEMENT type (regular | project)>
<ELEMENT prereq (course)>]
<ELEMENT takenby (student)>
```
The idea of attribute translation grammars (ATGs) is to treat the DTD as a grammar and recursively fire productions from the grammar to create an XML document. We now briefly review the syntax and semantics of ATGs (see [4]).

An ATG $\sigma : R \rightarrow D$ specifies a mapping from instances of the source relational schema $R$ to documents of the target DTD as follows: a) For each element type $A$ of $D$, $\sigma$ defines a semantic attribute $\$A$ whose value is a single relational tuple of a fixed arity and type; intuitively, $\$A$ controls the generation of $A$ elements in the XML view, and is used to pass data downward as the document is produced. b) For each production $P = A \rightarrow \alpha$ in $D$, $\sigma$ specifies a set of semantic rules, $\text{rule}(p)$. These rules specify the computation of the $B$ children of an $A$ element for each type $B$ in $\alpha$.

Given a database $I$ of $R$, the ATG $\sigma$ is evaluated top-down starting at the root $r$ of $D$. A partial tree $T$ is initialized with a single node of type $r$, and this node is marked unexpanded; we refer to unexpanded nodes as buds. The tree $T$ is then grown by repeatedly selecting a bud $b$ (of some element type $A$), evaluating the semantic rules associated with $A$, and marking $b$ expanded. Specifically, we find the production $P = A \rightarrow \alpha$ in $D$, and generate the children of $b$ by evaluating $\text{rule}(p)$ and using the value of the attribute $\$A$ of $b$. The rules $\text{rule}(p)$ are defined and evaluated based on the form of $\alpha$ as follows:

1. If $\alpha = B_1; \ldots; B_n$ then a node tagged $B_i$ is created for each $i \in [1,n]$ as a child of $b$. The tuple value of $\$B_i$ is associated with the new $B_i$ child is determined by projection from $\$A$. That is, $\$B_i = (\$A.a_1, \ldots, \$A.a_n)$ is in rule $p$ for $i \in [1,n]$, where $a_i$ is a field of the tuple $\$A$.

2. If $\alpha = B_1 + \ldots + B_n$ then $\text{rule}(p)$ is defined by

$$\text{case } f(\$A) \text{ of } \begin{cases} 1: (\$A.\text{null}, \ldots, \text{null}) \\ n: (\text{null}, \ldots, \text{null}, \$A) \end{cases}$$

where $f$ is a function that maps $\$A$ to natural numbers in $[1,n]$. That is, based on the conditional test, a node is created for exactly one child, $B_i$. The value of the parent attribute $\$A$ is passed down to that child. No $B_i$ child is created if $i \neq j$, and $\$B_j$ (the special value null) is ignored.

3. If $\alpha = B^*$, then $\text{rule}(p)$ is defined by $\$B \leftarrow Q(\$A)$, where $Q$ is an SQL query over a database of $R$, and it treats $\$A$ as a constant parameter. For each distinct tuple $t$ returned by $Q(\$A)$, a $B$ child is generated, carrying $f$ as the value of the $\$B$ attribute. To help ensure that only finite documents are created, only references to attributes and constants, but not expressions are allowed in the select-list of $Q$.

4. If $\alpha$ is $\text{str}$ then the rule specifies formatting of the values of $\$B$ for presentation (string/PCDATA). Such rules are not shown or discussed further.

5. If $\alpha$ is $\epsilon$, then no $\text{rule}(p)$ is defined and no action is taken.

The element children of the node $b$ become new buds and are also processed. The process proceeds until the partial tree cannot be further expanded: i.e., it has no unexpanded node. The fully expanded XML tree do not include attribute values $\$A$, which are only used to control the tree generation.

Example 2.2: The ATG $\sigma_0$ given in Fig. 3 defines the XML view described in Example 2.1. Here $\text{rule}(\text{course}).
db → course

$course ← Q_1$

$Q_1$: select distinct c.cno, c.title, 1 as tag from course c
where c.dept = 'CS'
union
select distinct p.cno, p.title, 2 as tag from project p
where p.dept = 'CS'

course → cno, title, type, prereq, takenBy

$@no = @course.cno, @title = @course.title,
@type = @course.tag,
@prereq = @course.cno,
@takenBy = @course.cno$

type → regular + project

$(\text{regular}, \text{project}) = \text{case type of}
1: (\text{type}, \text{null})
2: (\text{null}, \text{type})$

prereq → course

$@course ← Q_2(\text{prereq})$

$Q_2(c)$: select distinct c.cno, c.title, 1 as tag from prereq p, course c
where p.cno = c and p.cno ≠ c.cno

takenBy → student

$@student ← Q_3(\text{takenBy})$

$Q_3(c)$: select distinct s.ssn, s.name from enroll e, student s
where s.cno = c and e.ssn = s.ssn

Figure 3: Example ATG $\sigma_0$

rule(type) and rule(prereq) illustrate the cases (1), (2), (3) above. Given a registrar database, $\sigma_0$ computes an XML view as follows. It first generates the root element (with tag db), and then evaluates the query $Q_1$ to extract courses and projects of the CS department from the underlying database. For each distinct tuple $c$ in the output of $Q_1$, it generates a course child $v_c$ of db, which is a bud carrying $c$ as the value of its attribute $@course$. The subtree of the bud $v_c$ is then generated by using $c$. Specifically, it creates the cno, title, type, prereq and takenBy children of $v_c$, carrying c.cno, c.title, c.tag, c.cno and c.cno as their attributes respectively. It then proceeds to create a text node carrying c.cno as its PCDATA, as the child of the cno node; similarly for title. It determines the type of $v_c$ by examining c.tag; if it is 1 then a regular child of type is created; otherwise a project child is generated. It creates the children of the prereq node by evaluating the SQL query $Q_2$ to find prerequisites of the course, and again for each tuple in the output of $Q_2$ it generates a course node; similarly it constructs the takenBy subtree by evaluating $Q_3$ to extract student data. Note that $Q_2$ and $Q_3$ take c.cno as a constant parameter. Since course is recursively defined, the process proceeds until it reaches courses that do not have any prerequisites i.e., when $Q_2$ returns empty at the prereq children of those course nodes. That is ATG Handle recursion following a data-driven semantics. When the computation terminates the ATG generates an XML view as depicted in Fig. 2, which conforms to the DTD D0 given in Example 2.1.

As observed by [4], ATGS are more expressive than the view definition languages of previous publishing systems [11, 8].

Exception Handling. ATGS as defined refine the original definition introduced in [4]. An ATG of [4] may abort over a relational database, i.e., it may terminate unsuccessfully as the XML view it generates may violate the given DTD. In contrast, our revised ATGS do not abort. However, even our revised ATGS defined over a recursive DTD may not terminate. For example, $\sigma_0$ of Fig. 3 may not converge if the relation prereq in the underlying database is cyclic, e.g., if a course is its own prerequisite. Worse still, it is undecidable [4] to determine at compile time for an arbitrary ATG $\sigma : R \rightarrow D$, whether $\sigma$ terminates on all databases of $R$ [4]. To cope with this we introduce an exception handling mechanism. For an element type $A$ defined recursively, we consider a mild extension of its production $p = A \rightarrow a + \text{str}$. For example, we extend the production of course in the DTD D0 to be

course → (cno, title, type, prereq, takenBy) + str

In order to ensure termination, we modify tree generation to stop expanding the tree if a newly created node $v_c$ has the same type and semantic attribute value as one of its ancestors, $v$. In this case, we simply emit the string value of $v_c$’s attribute as the contents of the new node. It should be mentioned that this does not lose information, since the subtree of $v_c$, if constructed by following the original production $p = A \rightarrow a$, is just a copy of the subtree at $v$ and does not introduce any new information. This process is referred to as exception handling. Although exception handling slightly modifies the DTD embedded in an ATG, it ensures termination of ATG evaluation without loss of information.

Theorem 2.1: Let $\sigma : R \rightarrow D$ be an arbitrary ATG with exception handling. Then over all instances $I$ of $R$, $\sigma$ terminates and $\sigma(I)$ conforms to the DTD $D$.

Theorem 2.1 shows that an ATG $\sigma : R \rightarrow D$ is actually a total function: given a database $I$ of $R$, $\sigma(I)$ computes an XML document of $D$ that is unique up to reordering of $B$-elements produced by productions of the form $A \rightarrow B^*$. In the sequel we consider only ATGS with exception handling and refer to them simply as ATGS. Note that our incremental techniques also work on ATGS without exception handling.

3. External Trees

Incremental update makes sense only in the context of an externally maintained tree available to the client. A variety of potential implementations for the external tree exist, from a native storage system like Berkeley XML DB to an in-memory implementation of DOM. While in the current work we have used our own implementation of an in-memory tree in C++, we expect the data structures and algorithms we propose to apply to other implementations.

We next describe the external data structures maintained by our middleware to accept and process changes.

Node Identity. We assume that we can associate a compact, unique value with each tuple value taken on by a semantic attribute in $\sigma(I)$. We abstract away the implementation of this identity value by assuming, without loss of generality, the existence of a Skolem function gen_id (see, e.g., [9]) that, given the tuple value of a semantic attribute $\delta A$, computes id(A) that is unique among all identities associated with all semantic attributes (for example, it might encode the type and a unique value within that type).

Tree vs. Graph Representations. An important property of an ATG $\sigma : R \rightarrow D$ is that, for any database $I$ of $R$ and type $A$ of $D$, an $A$-element (subtree) $T_A$ in the XML view $\sigma(I)$ is uniquely determined by the value of the semantic attribute $\delta A$ at the root of $T_A$. Thus the ATG defines a function $st$ such that, given an element type $A$ and a value $t$ of $\delta A$, $st(A,t)$ returns a subtree rooted at a node tagged $A$ and carrying $t$ as its attribute.
Since a subtree $st(A \subseteq A)$ may appear at different places in the XML view $\sigma(I)$, if the middleware system is managing the external view, it may be more efficient to represent $\sigma(I)$ as a graph such that a single copy of $st(A \subseteq A)$ is stored and shared by its multiple occurrences. Indeed, if $\sigma(I)$ is stored as a tree, for example by using an available implementation of DOM [2], the tree may be exponentially larger than the graph representation. In light of this the current implementation of ATGS adopts the graph representation, but supports client navigation on the graph as if it were a tree. The trade-off of the space efficiency is that the support of tree navigation is complicated, as the path from a particular node to the root of the tree would be dependent on the route navigated; further, a mechanism must be provided to check for duplicates along this path to provide the exception handling semantics described above.

**Data Structures.** The external data structures used to represent the XML document are depicted in Fig. 4. The tree $T$ is stored in a hash index $H$ and a subtree pool $S$. Each entry of $H$ is of the form $(A, idA\_ptr)$, where $A$ is an element type, $idA$ is the unique id of a value of $A$, and $ptr$ is a pointer to the root node of the subtree $st(A \subseteq A)$ in $S$. The subtree pool $S$ consists of entries $(A, idA\_L)$, where $(A, idA)$ represents a node $v$ of $T$, and $L$ is a list $[(B_1, idB_1), \ldots, (B_n, idB_n)]$ representing the children of $v$ such that each $(B_i, idB_i)$ is an entry in $H$. Observe that in the graph representation, there is a one-to-one mapping from $H$ entries to the nodes in $S$.

**Handling Updates.** In this paper, we allow $\Delta I$ to be any group of updates to the underlying mass that preserves the consistency of the database (integrity constraints). XML updates $\Delta T$ generated from $\Delta I$ by one of the techniques described in the next two sections are represented as $(E^+, E^-)$, where $E^+$ is a set of edges to be inserted into the tree $T$, and $E^-$ is a set of edges to be deleted from $T$. The edges are represented as $(idA, idB, L)$, where $idA$ and $idB$ are the ids of the parent $A$-element and the child $B$-element respectively.

The system processes insertions as follows for each tuple $(idA, idB) \in E^+$: 1) find the $H$ entry $(A, idA\_L)$ and the pointer to the node $(A, idA\_L)$ in the subtree pool pool $S$; 2) insert $(B, idB\_L)$ into $L$ if it is not already in $L$, and 3) if $H$ does not have an entry $(B, idB\_L)$ in $H$, create an entry in $H$ and a node $(B, idB\_L)$ in $S$ with an empty child-list [ ]. Note that if there is already an entry $(B, idB\_L)$ in $H$, the new edge is actually a cross-edge to the existing $B$ node.

A subtle issue concerns the order of the children list $L$ when a $B$ node is inserted into $L$. Note that $L$ is constrained by the production $A \rightarrow \alpha$: if $\alpha$ is $B_1, \ldots, B_n$, the elements of $L$ are in the same order as their element types in the production; and if $\alpha$ is $B^\prime$, $L$ elements are ordered by a default order on the tuple values of the $B$ attributes of these elements.

The system carries out deletions similarly: for each tuple $(idA, idB) \in E^-$, 1) find the node $(A, idA\_L)$ in the subtree pool via the $H$-index entry $H(A, idA\_L)$, and 2) cuts the edge $(idA, idB)$ by removing $idB$ from the children list $L$. It should be noted that during the processing of relational updates $\Delta I$, no entries are physically deleted from $H$ because a) an entry may be shared by multiple subtrees of $T$, and b) as will be seen shortly, our incremental system minimizes unnecessary recomputations by reusing subtrees that have been disconnected from the tree by cuts.

**Garbage Collection.** Upon the completion of the computation of the new XML view $\sigma(I \oplus \Delta I)$, a garbage collection process runs in the background to remove $H$-entries and nodes in the subtree pool that are not linked to $\sigma(I \oplus \Delta I)$. Conceptually this can be done via a top-down traversal of the new XML tree, removing nodes that are no longer reachable from the root. In our implementation we associate with each $H$-entry a $count$ keeping track of the number of nodes linked to it, and maintain the counters when processing insertions and deletions; the garbage collection process removes unused nodes, i.e., those $H$-entries with $count = 0$ and their corresponding nodes in the subtree pool.

### 4. Pushing Incremental Work to the DBMS

In this section we present our reduction mechanism for incremental evaluation of ATGS. The approach is based on 1) a relational encoding of XML trees via a set of (interrelated) virtual attributes and edge relations and 2) a translation of an ATGS $\sigma : R \rightarrow D$ to a set $M_{\sigma_Q}$ of SQL 99 queries utilizing the with...recursive construct, which computes the attribute and edge relations of the XML views defined by $\sigma$.

Given this the reduction approach to incrementally maintaining an XML tree $T$ computed by ATGS $\sigma(I)$ works as follows: 1) encode with the attribute and edge relations, 2) map $\sigma$ to SQL 99 queries $M_{\sigma_Q}$, 3) define an incrementally updated materialized view in the source DBMS for each query of $M_{\sigma_Q}$, 4) in response to relational updates $\Delta I$, utilize the union functionality to capture the incremental changes made to each of the materialized views which can be directly transformed into the XML changes $\Delta T$ (i.e., $E^+$ and $E^-$), and 5) propagate the changes $\Delta T$ to the external tree as described in the previous section. This leads to a convenient approach to incremental evaluation of ATGS by pushing as much work as possible to the underlying DBMS.

We next focus on the relational encoding of XML trees and translation of ATGS to SQL 99 queries.

#### 4.1 A Relational Encoding of XML Trees

An XML view $\sigma(I)$ is encoded via node (semantic attribute) and edge relations.
Attribute Relations. The attribute relations are to capture the values taken on by the semantic attributes defined in σ(I). To avoid confusion, "attributes" of relational tables will be uniformly referred to as "columns".

Recall that σ associates with each element type B a tuple-structured variable $B$. For each such variable $B$, let gen$_B$ be a relation with columns matching the arity and type of $B$, along with a column for id$_B$ if an existing group of attributes does not serve this role. Further, in the context of a database instance $I$, assume that gen$_B$ is populated with all the (non-null) values taken on by $B$ during an evaluation of $σ$ on $I$.

We define each attribute relation, gen$_A$, in terms of a query $Q_{gen_A}$ involving the other attribute relations and the relations of $I$. To define $Q_{gen_A}$ we first rewrite sq1. queries embedded in $σ$ to queries that take gen$_A$ instead of a single tuple $A$, as a parameter. Specifically, consider productions $A → α$ in which $B$ appears on its right-hand side (rhs).

1. For productions of the form $A → B'$ with associated semantic rule $B \leftarrow Q(\Delta A)$, gen$_B$ is the union, over all values of $\Delta A$ in gen$_A$, of all $Q(\Delta A)$. In a manner similar to the level-at-a-time processing of [4], this query can be rewritten as $Q_i$, which takes no parameters but additionally accessors gen$_A$. This is accomplished, roughly by a) adding gen$_A$ to the from list of $Q$, and b) replacing references to $\Delta A$ in $Q$ with the corresponding references to gen$_A$.

2. If $B$ appears as some $B_i$ when $α$ is $B_1, \ldots, B_n$, then $Q'$ is a simple selection query from gen$_A$ that projects fields according to $f(\Delta A)$ (see Section 2).

3. If $B$ appears as some $B_i$ in $α$ is $B_1, +, B_2$, then $Q'$ can be written as select * from gen$_A$ where $f(α) = i$. Note that we assume further that $f(α)$ is computable in the dialect of sq1. used to express $Q_i$.

It is now simple to generalize this construction to handle the case where $B$ appears on the rhs of multiple rules in $σ$. Suppose that it appears on the rhs of rules $p_1, p_2, \ldots, p_m$, and that $Q_i$ is defined per the discussion above. Then gen$_A$ is formed by taking the (distinct) union of all $Q_i$ queries.

For example, Fig. 6 shows the definition of two attribute-relation generating queries when the resulting construction is applied to the $\alpha$ of Fig. 3.

We denote by AR$(σ)$ the set of all $Q_{gen_A}$ for $σ$.

Maintaining Edges. We next describe how to capture the edge relations of the xml view. Let edge$_{A\rightarrow B}$ be a relation with two columns id$_A$ and id$_B$. We create such a relation if $B$ appears on the rhs of some production for $A$. We overload the Skolem function gen$_id$ described in Sec. 3 to compute the unique id $α$ from a relational tuple $α$.

We now discuss how to derive queries to define these edge relations in terms of attribute relations and base relations. As before, consider first productions of the form $A → B'$, where $B' \leftarrow Q(\Delta A)$ is the associated semantic rule. In this case, edge$_{A\rightarrow B}$ is the set of pairs $(α, β)$ such that $α \in$ gen$_A$ with Q1(cmo.title,...) as (Q1 from Figure 3) with Q2(cmo.id, idp, p_idp, p_idc, srel) recursive as {
  select gen_id(cmo, c.title), 1) from Q1
union select gen_id(gp), gen_id(cmo, c.title), 1), null, null, 'course'
from Q2 gp, prereq p, course c
where p.cmno = gp.cno and p.cmno2 = c.cno
gp.srel = 'prereq'
union select null, null, cmo as pg_idc, cno as pg_idp, 'prereq'
from Q2 gc
where gc.srel = 'course'
select distinct cmo, title, tag
from Q2, course
where Q2.srel = 'course'
and Q2.idc-course.cmo
}

with Q1(cmo.title,...) as (Q1 from Figure 3)
with Q2(cmo.id, idp, p_idp, p_idc, srel)
recursive as {
  select gen_id(cmo, c.title, 1) from Q1
union select gen_id(gp), gen_id(cmo, c.title, 1),
null, null, 'course'
from Q2 gp, prereq p, course c
where p.cmno = gp.cno and p.cmno2 = c.cno
and gp.srel = 'prereq'
union select null, null, cmo as pg_idc, cno as pg_idp, 'prereq'
from Q2 gc
where gc.srel = 'course'
select distinct cmo, title, tag
from Q2, course
where Q2.srel = 'course'
and Q2.idc-course.cmo
}

4.2 Computing Attribute and Edge Relations with sq1. 99 Linear Recursion

As the running example illustrates, the attribute relations are potentially mutually recursive, and the edge relations depend on the attribute relations. Furthermore, each attribute relation can be related to itself and other attribute relations through a variety of paths, raising the possibility that concluding evaluation of updates to the attribute relations will be excessively complex. Fortunately, this is not the case.

While the attribute relations are recursive there are substantial constraints on this recursion. Consider the references in the query $Q_{gen_A}$ to other attribute relations $B, C, \ldots$. First, observe that such a relationship cannot be made via negation (in the Datalog sense), that is the reference to $B$ cannot appear in a not exists clause in $Q_{gen_A}$. Second, note that references to two different attribute relations for example $B$ and $C$, must be made in two different subqueries of $Q_{gen_A}$ such that these subqueries are joined only by the top-level union operator introduced by the construction. Finally since expressions are not allowed in the select clauses of view-definition queries, aggregate expressions which depend on attribute relations will not appear.

Given these restrictions, the computation of the attribute and/or edge relations can be accomplished with the sq1. 99 with..recursiv construct. Let $Q_0$ be a set of edge and attribute relation-generating queries from $σ$ which either includes all of the attribute relations or all of the edge relations. Let $G_0$ be a graph which in which there is a node for each query in $Q_0$, and an edge from $Q_0$ to $Q_1$ iff $Q_0$ refers to the virtual relation gen$_A$ in the from clause. Let $c_0, \ldots, c_n$ be the connected components of this graph. We now merge components and generate possibly recursive queries along the same lines as [17]. First, single component with only a single incoming edge is merged with the source component. Let the result of this process be $c_1, \ldots, c_m$. For each such component $c_i$, if it is acyclic, then a set of merge-queries
5. Bud-Cut Incremental Evaluation

As discussed in the introduction, the reduction approach is not practical for middleware-based XML publishing since it depends on a combination of features not yet found in even the most advanced commercial DBMSs. While middleware should depend on only the most common features. This observation motivates us to propose the Bud-Cut approach to incrementally evaluating $\text{ATG} \sigma : R \rightarrow D$, which does not require the underlying DBMS to support with...recursive.

A middleware system based on the bud-cut mechanism is depicted in Fig. 8. The system interacts with an underlying DBMS and maintains a hash index $H$ and a subtree pool for the external XML view $T(\sigma(I))$ as described in Sect. 3. It responds to a relational update $\Delta I$ in three phases. The first phase, bud-cut generation, identifies the portions of the existing tree that will be affected by the updates $\Delta I$, and propagates $\Delta I$ to XML changes to the existing nodes in $T$. Nodes created in this phase are "buds", i.e., they are marked unexpanded; and nodes to be removed are "cuts". The second phase, bud completion, constructs subtrees under buds (and generates new buds), taking advantage of ATG properties to avoid recompilation and reusing existing subtrees when possible. The third phase, garbage collection, runs after bud completion is finished and removes unreachable subtrees.

While requiring several round-trips between the middleware and the DBMS, this approach is able to exploit other optimizations by taking advantage of the specifics of ATG semantics and XML views. In particular, as observed in Sect. 3, the value of a subtree $\text{srt}(A, t)$ is determined by the tuple value $t$ of the semantic attribute of the subtree's root node, changes to the children of the existing nodes in the tree can be computed by a fixed set of non-recursive queries. Furthermore, any new subtree $\text{srt}(A, t)$ can be reused and thus needs to be computed at most once. In addition, the bud-cut approach allows the update process of the external tree $T$ to run in parallel with the compilation of XML changes $\Delta T$, thus improving the response time. This last point leads to a variety of options for lazy evaluation of incremental updates. That is, the middleware can optionally defer complete processing of updates until the subtrees affected by these updates are accessed. When the external view is materialized as an in-memory tree, it is possible to support tree navigation concurrent with processing.

5.1 Implementation of Delta SQL Queries

Assuming the existence of virtual attribute relations, the first step is to derive two sets of incremental (nonrecursive) sq. queries from edge-generating queries $ER(\sigma)$. These two forms of incremental queries will be used in the bud-cut generation phase and the bud-completion phase, respectively.

Incremental queries of the first form are derived as follows. For each query $Q_{\text{edge recount}}$ in $ER(\sigma)$, we define a bud generating query formed by adding all the columns of the attribute $\$B$ to the select list of $Q_{\text{edge recount}}$. An incremental form of this bud-generating query, $\Delta Q_{\text{edge recount}}$, is then created by using a counting method like [15]. In a nutshell, [15] associates a count with each tuple in a view to keep track of the number of alternative derivations of the tuple via the view, and computes updates to the view, namely, insertions and deletions of tuples, by incrementing or decrementing the counts of its tuples in response to changes to base relations.

Since the method of [15] is assumed to execute in the DBMS, it assumes access to both $I$ and $I \oplus \Delta I$. However, our middleware system is separate from the DBMS and can only access $I \oplus \Delta I$. To find $I$, for a given relation, $R_i$, we assume

\[
\begin{align*}
\text{create view oldcourse as} \quad & \text{select * from course} \\
& \text{where not exists ( select 0 from \Delta course dc} \\
& \text{where dc.cno = course.cno ));} \\
\Delta Q_{\text{edge recount}}: \quad & \text{select p.cnt, gen_id(gp), gen_id(c.cno, c.title, 1), c.cno, c.title, 1 as tag} \\
& \text{from gen_preReq gp, \Delta course p, course c} \\
& \text{where p.cno1 = gp.cno and p.cno2 = c.cno} \\
& \text{union} \\
& \text{select c.cnt, gen_id(gp), gen_id(c.cno, c.title, 1), c.cno, c.title, 1 as tag} \\
& \text{from gen_preReq gp, old_preReq p, \Delta course c} \\
& \text{where p.cno1 = gp.cno and p.cno2 = c.cno}
\end{align*}
\]
the existence of a change table $\Delta I$, that holds $\Delta I$ restricted to $R_i$ as well as a count cnt with the value either "$+$" or "$-\$" to indicate an insert or delete. Further, we assume that $\Delta I$ has already been applied to $I_i$, the instance of $R_i$ in $I$.

In order to simulate the pre-image of $I_i$, we define a view old$R_i$: for each relation $R_i \in R$, which simply selects tuples of $I_i$ that do not appear in $\Delta I$. It can be implemented efficiently if $\Delta I$ is small and is indexed on the key of $R_i$.

For example, Fig. 9 gives the incremental bud-generating query for edge$\_prereq$ course, which is derived from the query Qedge$\_prereqcourse$ by using the counting method of [15]. It computes updates to edge$\_prereq$ course in response to changes Acourse and $\Delta$prereq. Observe that in addition to edges of edge$\_A$,$B$, the query also returns the corresponding $SB$ values. A query $\Delta Q_{edge\_A}$,$B$(Bu$d_a$) to compute changes to edge$\_A$,$B$ in response to a set Bu$d_A$ of insertions to the attribute relation gen$\_A$. The query is derived from $Q_{edge\_A}$,$B$ by substituting Bu$d_a$ for gen$\_A$ and by adding all the columns of the attribute $SB$ to the select list of $Q_{edge\_A}$,$B$. It computes the edges of the newly inserted $A$-nodes in Bu$d_a$. In contrast to the incremental bud-generating query $Q_{edge\_A}$,$B$, the incremental query $\Delta Q_{edge\_A}$,$B$(Bu$d_a$) takes Bu$d_A$ as a parameter and does not access gen$\_A$; furthermore, it uses the new relation $I \ominus \Delta I$ without requiring the old $I$. For example, Fig. 10 shows the incremental bud-completing query of $Q_{edge\_completing}\_BU$ which assumes that course has been updated by $\Delta$course.

5.2 Bud-Cut Generation

Having derived the incrementalized SQL queries, this phase proceeds as follows:

1. The incremental bud-generating queries $\Delta Q_{edge\_A}$,$B$ are executed in the DBMS, yielding changes to the edge relations a set of insertions $E^+$ and deletions $E^-$. The sets $(E^+, E^-)$ are transmitted to the middleware and posted against the tree $T$, as described in Sect. 3. Moreover, the $SA$-value of each new $A$-node added to $T$ is also added to Bu$d_a$, a set of buds of type $A$. An $A$-node is marked unexpanded if its $SA$ value is in the Bu$d_A$ set.

2. The bud-cut approach does not require the attribute relations to be materialized in the DBMS since they can be computed from the hash table and the subtree pool in the middleware. To speed up the response time with a tradeoff of space, we treat these relations as views maintained by the DBMS. In the latter case, the changes computed in step (1) above also update these views; furthermore, in step (2) tuples in Bu$d_a$ are sent to the DBMS and inserted into gen$\_A$ for all element types $A$ in the ATG.

It should be remarked that the bud-cut-generation phase updates the parent-child edges of all the existing nodes in the XML view $T$ in response to relational changes $\Delta I$. Thus all the edge deletions are handled in this phase, by evaluating the incrementalized SQL (bud-generating) queries of $ER(\sigma)$ once. Note further that nodes disconnected by the edge cuts are not removed from $H$-index or the tree but instead remain in the subtree pool pending reuse during bud generation. Garbage collection of the nodes that have been disconnected from the tree will be handled by a background process after the bud-completion phase, as described in Sect. 3.

A complicated STDM may lead to a large number of incremental bud-generating queries. For an given small relational change $\Delta I$, however, one only needs to evaluate those that refer to a relation affected by $\Delta I$. Techniques [21, 6] for identifying queries irrelevant to $\Delta I$ can further reduce the number of bud-generating queries that need to be evaluated.

Example 5.1: Recall the ATG $\sigma_0$ defined in Fig. 3. Assume that an XML tree $\sigma_0(I)$ of a registrar database $I$ is maintained by the middleware. Consider relational changes $\Delta I$: deletion of a CS course from the course relation, along with deletions from prereq and enroll accordingly. Given this the bud-cut-generation phase computes only $E^-$, which consists of deletions of edges (db.course), (course.prereq), etc. These changes are made to the XML tree $\sigma_0(I)$ in this phase.

Note that no buds are generated, i.e., Bu$d_a$ is empty for all $B$ in the ATG. In other words, although the relational changes have impact to the XML tree at an arbitrary level, they are captured in the bud-cut-generation phase by evaluating a fixed number of incrementalized SQL queries.

On the other hand, if $\Delta I$ is to update the name fields of some student tuples, then it involves both deletions and insertions, i.e., Bu$d_{student}$ is nonempty and consists of student buds. The bud-cut-generation phase handles deletions and the next phase the bud-completion phase, proceeds to generate subtrees of new student buds.

5.3 Bud Completion

The previous bud-cut-generation phase creates sets Bu$d_a$ consisting of the $SA$-attribute values of new $A$-nodes for each element type $A$ in the ATG, and the bud-completion phase is to produce the subtrees of these buds. The process may further generate new buds, but does not incur deletions.

The processing of bud completion is conducted by Algorithm eval, given in Fig. 11. The algorithm takes Bu$d_a$'s from the generation phase as input, and processes each non-empty Bu$d_a$ based on the production $p = A \rightarrow a$ and its associated rel$\_p$ (cases 1-4), generating children of these $A$-elements. For example, in case 3, it processes a given set Bu$d_a$ as follows. It first finds the edges of the $A$ nodes in Bu$d_A$ that are to be inserted into edge$\_A$,$B$ by evaluating the incremental bud-completing query $\Delta Q_{edge\_completing}\_BU(Bu$d_A$) in the DBMS. It then updates the children lists of these $A$ nodes in the subtree pool $T$, taking advantage of the $H$-index and id$\_A$. Then for each $B$ child id$\_B$ of such an $A$-bud id$\_A$, it invokes the procedure process$\_B$ (id$\_B$) to inspect whether there already exists an entry for the node (B, id$\_B$) in the hash index $H$. If so, it simply adds a cross edge (id$\_A$, id$\_B$) instead of recomputing the subtree of (B, id$\_B$); otherwise, it creates a new $H$-entry for (B, id$\_B$) and adds the attribute value $SB$ of id$\_B$ to Bu$d_B$. This yields Bu$d_B$, the set of new $B$-buds that will be further expanded at the next level of the tree. Note that to populate Bu$d_B$ we need the attribute values $SB$ of these newly inserted $B$ nodes.

Observe the following properties of Algorithm eval.
\textbf{Input:} newly inserted nodes $Bud_A$ for all element types $A$ in $\sigma$.

\textbf{Output:} completed subtree for all nodes in $Bud_A$.

1. repeat
   2. for all nonempty $Bud_A$ with $A$ in $\sigma$ do
   3. case the production $p : A \rightarrow \alpha$ of
   4. \hspace{1em} (1) $A \rightarrow B_1, \ldots, B_n$;
   5. \hspace{1em} for each tuple $t$ in $Bud_A$ do
   6. \hspace{2em} $S_B \leftarrow f(t);$ /* $f$ in rule(p) for $S_B$ */
   7. \hspace{2em} \textbf{process} $(B_1, S_B)$;
   8. \hspace{2em} let $v = H(A, t)$;
   9. \hspace{2em} if $v \notin H(A, t)$ then $v = H(A, t)$;
   10. \hspace{2em} \textbf{end process} $(B_1, S_B)$;
   11. \hspace{1em} \hspace{1em} $v = (B_1, t)$;
   12. \hspace{1em} $v.L = \{B_1, t\}$;
   13. \hspace{1em} \textbf{end for}
   14. \hspace{1em} \textbf{end case}
   15. \hspace{1em} \textbf{end for}
   16. \hspace{1em} \textbf{return}

\textbf{Procedure} \textbf{process} $(A, t)$.

\textbf{Input:} element type $A$, tuple $t$ as a value of $S.A$.

\textbf{Output:} modified index $H$, subpool tree and the sets $Bud_A$.

1. \hspace{1em} $tid := \text{gen.id}(t);$ /* $\text{gen.id}$: Skolem function */
2. \hspace{1em} if $(A, tid)$ is an entry of $H$
3. \hspace{1em} then return;
4. \hspace{1em} create a node $v = (A, tid [ ];$ in the subpool tree;
5. \hspace{1em} create an entry $(A, tid)$ in $H$;
6. \hspace{1em} $H(A, tid).pdr = &v;$ /* $\text{address of } v$ */
7. \hspace{1em} $H(A, tid).L = [];$ /* $\text{empty list}$ */
8. \hspace{1em} $Bud_A := Bud_A \cup \{t\}$; /* newly created nodes */
9. \hspace{1em} \textbf{return};

\textbf{Figure 11: Algorithm eval}

(1) At each iteration (level), any paths in the XML view that begin at the root and do not encounter a \texttt{bud} are guaranteed to be correct, and thus partial results of the new XML view $\sigma(I \oplus \Delta I)$ can be exposed to the users in parallel with the computation of $\Delta T$.

(2) It minimizes unnecessary computations via the procedure \textbf{process}; it reuses subtrees that have been computed either by Algorithm eval or earlier for the old view $\sigma(I)$. Thus each bud is computed at most once. The reuse of the previous computations is possible since no nodes are removed from the $H$-index or subpool tree at this stage.

(3) The bud-completion phase does not need materialization of the attribute relations, since all the $S.A$-values that $\Delta Qedge_{bud, A} \, B$ needs are in $Bud_A$, not in $gen_A$. Note that only case (3) in Algorithm eval needs to access the DBMS.

(4) The procedure \textbf{process} also ensures that algorithm eval always terminates, no matter whether the input \texttt{ATG} has exception handling or not. Specifically, if a descendant \texttt{bud} of an $A$-node $v$ with attribute value $\$A$ is also an $A$ node with the same $\$A$ value, the node \texttt{bud} is not created since there is already an entry for $(A, \$A)$ in the $H$-index; instead, any edge to \texttt{bud} is treated as a cross edge and linked to $v$. Thus no infinite computation is incurred.

Example 5.2: Consider again the $\texttt{ATG}$ $\sigma A$ and relational changes $\Delta I$ that update the name fields of some student tuples. The bud-cut-generation phase generates a set $Bud_{budgen}$ consisting of these updated $\texttt{student}$ tuples, generates new $\texttt{student}$ buds for these tuples, and redirects affected edges to these new nodes. The bud-completion phase completes the subtree of these new $\texttt{student}$s by creating new name fields and reusing sax fields for the students. This is done in one iteration of the outer loop of Algorithm eval, although the $\texttt{student}$s may appear at arbitrary levels of the XML view.

Now consider changes $\Delta I$ consisting of insertion of a new CS course to the course relation, along with insertions of its prerequisites to $\texttt{prereq}$ and insertions to enroll for the students who have enrolled in the course. The bud-cut-generation phase generates a set $Bud_{budgen}$ consisting of the newly inserted course tuple, creates a new node $c$ representing the course and adds an edge from the root $\texttt{db}$ to $c$. The bud-completion phase takes $Bud_{budgen}$ as input and constructs the subtree of $c$. Since all the prerequisites of $c$ are already in the XML view, the subtree of $c$ can be completed in two iterations of the outer loop of Algorithm eval. In the first iteration only $Bud_{budgen}$ is nonempty, and the algorithm creates new $\texttt{cnr}$, $\texttt{title}$, $\texttt{prereq}$, $\texttt{type}$ and $\texttt{enrolled}$ children of $c$. The second iteration completes their subtrees by reusing the existing course and student nodes. Since no new nodes are generated in the second step, the iteration terminates although the subtree under $c$ may have a depth greater than two due to its prerequisites hierarchy.

\textbf{5.4 Implementation Issues}

\textbf{Overlapping Phases.} While the first and second phase have been presented as completely separate, in practice it may be advantageous to "pipeline" them. In particular, consider the case where $\Delta I$ affects base tables for both $\Delta Qedge_{a,b} \cup \Delta Qedge_{a,b,c}$ and $\Delta Qedge_{b,c} \cup \Delta Qedge_{b,c,d}$. It makes sense to compute $\Delta Qedge_{a,b,c}$ first, allowing the $A$ nodes in the tree to be expanded and thus $\Delta Qedge_{b,c}$ to be incremented before $\Delta Qedge_{b,c,d}$ is executed. This follows the observation in [15] that incremental queries should be evaluated according to their "stratum number", which reflects dependency relationship.

\textbf{Query Merging.} Query merging has been proved useful for systems to publish relational data in XML [11, 4]. The idea is to reduce the number of queries issued to the DBMS by merging multiple queries into a single larger query via outer-join/outer-union. Query merging can help decrease the communication costs between the middleware and DBMS, while also potentially diminishing query processing time and execution overheads.

The generation of subtrees under buds is similar to the original process of ATG generation [4], and query merging can be used in the bud-completion phase to merge multiple edge-generating queries of the form $\Delta Qedge_{bud, A} \cup \Delta Qedge_{bud, B}$ into one. However, the reuse of existing subtrees is an important difference, and direct application of the merging techniques for ATG [4] may lead to unex-
6. Experimental Results

Our experimental evaluation focuses on the effect of database and update size on the performance of our approach and on the effectiveness of the proposed subtree caching strategy.

We build the source database based on the schema of Example 2.2, but with an additional *taughty* table giving instances of professors teaching courses. The database size in terms of the number of courses, with scaling factors for other tables given in Fig. 12. A smaller pool of $f_{ij}$ of the courses are used to initially build the tree. In effect, $f$ controls the probability a new prerequisite will exist in the tree already. For the recursive part of the schema, between one and three random prerequisites with lower costs from the pool are generated for each course. The maximum depth of this recursive part is limited to 8 levels of XML nodes.

An update consists of inserting or deleting a course and its associated prerequisites. Fixed size batches of meaningful updates, $w = \Delta I$, are generated by ensuring that inserts are not duplicates and that deleted courses are present. Since deletes in our system generally execute much faster than inserts, $w$ is constrained to have 50% inserts unless stated otherwise. The locality of updates is controlled by selecting the updates in $w$ from a random pre-generated "universe," $S$, of courses. When $S$ is smaller, the updates in $w$ are more related. The parameters and their default values are summarized in Fig. 13. Experiments are performed by running sequences of such batches for two workloads $W_1$ and $W_2$, and using *bud-cut* to propagate the changes to the external tree. The first workload $W_1$ consists of executing such batched updates against our example ATG, $\sigma_0$. A second workload $W_2$ modifies $\sigma_0$ to restrict prerequisites to courses taught by instructors from the computer science department. Thus, $W_2$ introduces a join to *taughty* in the recursive part of the view.

The experiments were run on a system with a 2GHz Pentium IV processor with 1G bytes of RAM; they were conducted with a large (256MB) *warm* DMBS buffer cache. While incremental update may reasonably avoid disk access since the relevant data was just updated, fully cached operation for document publishing is also reasonable, at least for published XML data up to a few tens of megabytes in size given modern database configurations. Similarly, the external XML tree (see Sect. 3) is fully cached in RAM on the same system. Each experiment was run five times and the average

### Table 1: Table sizes

<table>
<thead>
<tr>
<th>course</th>
<th>cardinality</th>
<th>tuple size</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>40 bytes</td>
<td></td>
</tr>
<tr>
<td>1-3/coues</td>
<td>8 bytes</td>
<td></td>
</tr>
<tr>
<td>student</td>
<td>100</td>
<td></td>
</tr>
<tr>
<td>takenby</td>
<td>5-10/courses</td>
<td>8 bytes</td>
</tr>
<tr>
<td>taughty</td>
<td>1/2 courses</td>
<td>8 bytes</td>
</tr>
</tbody>
</table>

### Table 2: Table of symbols

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Meaning</th>
<th>Default val.</th>
</tr>
</thead>
<tbody>
<tr>
<td>$S$</td>
<td>the &quot;universe&quot; of the updated courses</td>
<td>300</td>
</tr>
<tr>
<td>$f$</td>
<td>the percentage of $S$ in the published XML view</td>
<td>50%</td>
</tr>
<tr>
<td>$w$</td>
<td>a unit of work containing both insertions and deletions</td>
<td>[w] = 100</td>
</tr>
<tr>
<td>$I$</td>
<td>the database size</td>
<td>$</td>
</tr>
</tbody>
</table>

is reported here. All numbers obtained for each average are within 6% of the average value. Unless otherwise noted, we used the proposed caching strategy.

We compare the cost of incremental update against the cost of full tree recreation for a variety of database sizes. We scale the database by varying $|I|$ from 100 to 10K, and setting $[w]$ to 4% of $|I|$. For this experiment, locality is not considered ($S = I$, $f = 100%$). Figure 14(a) shows the resulting time with $|I|$ on a log-scale for workloads $W_1$ and $W_2$. Not surprisingly, incremental update far outperforms full reconstruction for small updates once the overall database size gets appreciably large. Note that incremental update for $W_2$ is faster than for $W_1$ since the tree is significantly smaller (and thus less update work in the tree is needed).

Fixing $|I| = 10K$, we vary $|w|$ and study the behavior of incremental update vs. full tree recreation. The result is presented in Fig. 14 (b). As this figure shows, for $W_1$ and $W_2$, *bud-cut* scales nicely with the size of the change in the source database, $|w|$. We note that the *bud-cut* approach is sensitive to the selectivity of the ATG queries, and low selectivity leads to substantial activity on the hash index in the subtree pool. Accordingly, the incremental update and the full tree recreation for $W_1$ and $W_2$ cross when 10% and 80% of the database is updated, respectively.

To investigate the impact of the subtree pool, we evaluate the performance impact of turning it off. We fix $|w| = 100$ and vary $|S|$ from 100 to 600 (decreasing the locality of updates as $S$ grows) for $W_1$. (In this experiment, the percentage of inserts in $w$ is allowed to vary.) Finally, we consider two values of $f$, 50% and 80%, to control the probability that a given subtree will already appear in the tree. The results are shown in Fig. 15(a), where each point represents the average of 20 such experiments, with garbage collection fired after each run. As expected, the impact of subtree caching is greatest with smaller values for $|S|$ and $f$. With $f = 50\%$, the impact is substantial across the range. However, the curves for $f = 80\%$ are rather flat, since most cached subtrees are small due to "natural caching" by the rest of the tree.

While in Fig 14(b) we compare the size of the change in the source database to the time taken, it is also interesting to consider the size of the change in the *output* XML tree and compare this to the time taken for the incremental update. An experiment to investigate this relationship is shown in Fig. 15(b) (for workload W1). To capture this we output the XML tree before and after the incremental update $w$. We then use X-Diff [32] to produce an edit script converting the old tree to the new tree and measure its size as an approximation of $|\Delta T|$. Furthermore, $|w|$ remains 100, $f$ is 50% and $|S|$ is set to 300. This experiment shows that *bud-cut* also behaves reasonably well with respect to output size, and further, that the *subtree caching* optimization (in combination with the graph-oriented storage model) tends to be more effective exactly when an incremental update to
the database causes a large change in the output tree.

7. Related Work

Incremental computation has proved useful in many areas (see [27] for a survey). In particular, incremental view maintenance algorithms have been extensively studied for relational, datalog (see [13] for a collection of readings) and object-oriented views (see, e.g., [12, 18]). On the one hand, those algorithms are not directly applicable to incremental evaluation of XML publishing. Unlike traditional database views, XML views (ATGs) are defined by associating a collection of SQL queries with a (possibly recursive) DTD. These views are stored outside of DBMS, and their incremental maintenance is to propagate relational group updates to external XML trees. On the other hand, our incremental evaluation algorithms leverages previous work for traditional databases. Our reduction approach relies on the support of incremental maintenance of sqi.9 views by the underlying DBMS, and our top-down mechanism makes use of the counting algorithm of [15] to incrementally evaluate sqi.9 queries.

We now draw the analogy of our top-down ATG evaluation algorithm to incremental evaluation of recursive queries. A number of algorithms have been developed for evaluating recursive datalog queries and prolog programs notably [15, 29]. These algorithms typically involve a first phase where conservative deletions are conducted, followed by a second phase where insertions are performed, which may restore information deleted in the first phase. In contrast, our algorithm avoids recomputation between the two phases to an extent by reusing subtrees computed earlier and by deferring removal of disconnected subtrees. For example consider an extreme case where the conservative deletion involves deletion of the root of the tree, and thus the entire tree, whereas in the second phase the tree is reassembled with somewhat differing children of the root. While this may require reconstruction of the entire tree for the previous algorithms [15], our algorithm is capable of making maximum reuse of existing subtrees without doing unnecessary recomputations. Another difference between our algorithms and previous ones is that the conservative deletion phase repeatedly executes deletion operations for an unbounded number of times [15, 29], whereas by taking advantage of the semantics of XML trees, our bud-cut-generation phase does all the deletions by evaluating a fixed number of incrementalized SQL queries.

It is worth remarking that the possibility of encoding ATGs with sqi.99 queries was first suggested by [17], but [17] did not show how the encoding should be defined. It should be mentioned that there is a big gap between XML trees computed by an ATG and their relational encoding, while the former can be exponentially large in the size of the relational database. The latter is bounded by a polynomial in the database size. Thus an sqi.99 encoding of an ATG alone does not suffice to compute XML trees defined via ATG.

There have also been incremental maintenance algorithms for semistructured views [1, 31, 33]. These algorithms are developed for views defined with a nonrecursive query over graph structures, and cannot be applied to ATG evaluation. Incremental update of external views has been studied for data warehousing [22, 30, 34] and web site maintenance [19, 20]. Techniques developed in this work are complementary to ours, addressing e.g., compensating queries to cope with updates to decoupled sources, maintaining mediated views via constrained rules, scheduling updates for a group of views to maximize data quality, and deciding where to materialize web views with respect to access and update patterns.

There has also been a host of work on the self-maintainability of relational and datalog views (e.g., [21, 6]). We plan to study techniques for detecting updates irrelevant to XML views in order to optimize XML view maintenance.
The only work on maintaining XML views that we are aware of is [10]. The views considered in [10] are defined via a simple nonrecursive algebraic query over XML trees and as a result, maintenance of the views can be done via a bottom-up traversal of the source XML tree. These views are not capable of expressing ATGs, and the techniques of [10] cannot be applied to incremental evaluation of XML publishing of relational data. There has also been recent work on incremental XML validation [26], which differs from our work in that it focuses on validating XML documents in response to a single insertion or deletion of XML subtrees, rather than propagation of relational (group) updates to XML views.

Finally, the strategy commonly used by XML publishing middleware that pushes simple queries to the DBMS and conducts the rest of the work for composing a view at the middleware traces back to query processing for distributed relational databases [5]. Again recursive XML views introduce new challenges not encountered in the relational context.

8. Conclusion

We have proposed two approaches for incremental evaluation of ATGs [4]. The reduction approach is based on a non-trivial translation of ATGs to SQL 99 queries with recursion. Upon the availability of the support for incremental maintenance of SQL 99 views by commercial DBMS, the reduction approach leads to a convenient mechanism to maintain external XML views without requiring sophisticated middleware implementation. In the absence of high-end DBMS features, the bud-cut approach provides novel algorithms and a middleware system for efficient incremental maintenance of XML views. The algorithms and system minimize unnecessary recomputations by capitalizing on the ATG semantics and optimization techniques developed for XML publishing. We have implemented the bud-cut middleware, and our experimental results demonstrate that our approach, algorithms and optimization techniques are effective for maintaining XML views. To the best of our knowledge, our work yields the first effective framework for incremental evaluation of schema-directed XML publishing of relational data.

We plan to extend the current work in a number of directions. First, we note that our assumption of an in-memory hash table limits the technique for extremely large documents cached in middleware. We plan to address this by developing a streaming version of the incremental update so that an existing document can be read from disk and updated in a single pass. Second, the more complicated XML Schema standard is gaining popularity, and our next step toward handling XML Schema-directed publishing will be to extend our bud-cut algorithm to accommodate XML integrity constraints. Finally, we are planning to extend the present algorithms for use in schema-directed XML integration [5].

9. References